

#### the key value index optimized for size and speed



#### based on finite state (FST, immutable data structure)

Opensource (Apache 2.0)

written in C++(core), Python(binding)

# keyvi @CLIQZ Search in the Browser







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## keyvi in numbers @CLIQZ

1.8bn URLs

13bn data points (key value pairs)

2.5TB index size

60ms average latency

6k requests per minute

99.9% availability







#### $(\rightarrow 10/2014)$ to slow/heavy no need for full-text search did not fit our model





(→ 10/2014) to slow/heavy no need for full-text search did not fit our model



(4/2014 →) simple space efficient server side scripting





(→ 09/2015) capacity problems single threaded heap based single-point of failure





keyvi



(→ 09/2015) capacity problems single threaded heap based single-point of failure



 $(05/2015 \rightarrow)$ massive size reduction multi threaded/process shared memory more reliable



#### Size Comparison

# k/v pairs in million Redis keyvi 1 456 385 10 4538 3742 100 45303 36327

\* Size in MB



# **Compression Ratio per Type**





#### Lookup Benchmark 200 100 Chunksize Redis keyvi (RPC, 1 thread) 50 do your own benchmarkel 10 0 100000 200000 300000 400000 500000 600000 Requests per second client/server on the same machine chunksize == size of (redis) pipeline host type: r3.8xlarge(AWS)



## Scaling with Redis





www.keyvi.org

# Scaling with Redis 1 Machine



every Redis process has its own heap data access is local

if Redis dies, data has to be reloaded (can take a significant amount of time)



# Scaling with keyvi



keyvi is shared memory

several processes/threads can read

a crash effects 1 process

no deserialization/loading

no need for IPC/networking if local



# keyvi on SSD









#### Finite State in keyvi

#### Under the Hood: FST

#### An FST consisting of: berlin, buzzwords, buzz, keywords





### Under the Hood: a Trie for Comparison

An FST consisting of: berlin, buzzwords, buzz, keywords





#### **FST Incremental Construction**

does not fit in memory

- -> we need all outgoing transitions before persisting
  - -> (external) sort upfront

minimize on the fly

-> hashtable with low overhead, bounded, apply LRU

stream input and output



#### **Compact Persistence**

store nodes as compact as possible

-> use clever bit magic

Lucene: list of small byte representations

Keyvi: sparse array packing



#### Sparse Array









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Α	В		D	F			K	L	Μ	Ζ			R	Т			
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#### Sparse Array Tricks

fast bit vector comparison (intrinsics)

variable length encoding

relative offsets

sliding windows

. . .



# Keyvi Lookup

#### exact match == n \* access in both arrays





offset in a separate buffer stored as part of the finite state supported values: key-only, ints, strings, json special subtype: int with inner weights json store: msgpack, compression (snappy, zlib)



#### **Updates and Realtime**

#### (Linus Torvalds, 1998)

Almost all "real time" people end up actually being perfectly ok with "fast enough", and very very few people are really \_hard\_ real time.

#### Simple Updates @CLIQZ

Delta Update t-2

Delta Update t-1

Delta Update t (as needed: hourly, daily, on-demand)

Main Index / Master Segment (monthly)

under development: keyvi merger



exact matching entity recognition, e.g. in pySpark

approximate matching: close/near match e.g. for Geo applications scoring based: Levenshtein & Co

completion matching: prefix, multi-word, approximate





#### More infos/material/code at http://keyvi.org



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#### **Benchmark FST**



